Rigorous numerical computation of polynomial differential equations over unbounded domains

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Joint work with O. Bournez and D. Graça

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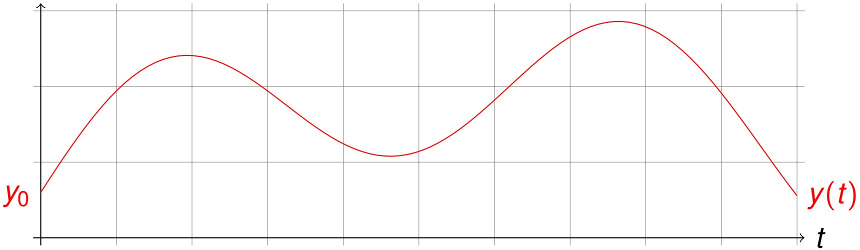
Outline

Ordinary Differential Equations (ODEs)



System of ODEs:

$$\begin{cases} y_1(0) = y_{0,1} \\ \vdots \\ y_n(0) = y_{0,n} \end{cases} \qquad \begin{cases} y_1'(t) = f_1(y_1(t), \ldots, y_n(t)) \\ \vdots \\ y_n'(t) = f_n(y_1(t), \ldots, y_n(t)) \end{cases}$$

More compactly:

$$y(0) = y_0 \qquad y'(t) = f(y(t))$$

Computability

Let $I = [0, a[$ and $f \in C^0(\mathbb{R}^n)$. Assume $y \in C^1(I, \mathbb{R}^d)$ satisfies $\forall t \in I$:

$$y(0) = 0, \qquad y'(t) = f(y(t)). \tag{1}$$

Can we compute $y(t) \pm 2^{-n}$ for all $t \in I$ and $n \in \mathbb{N}$?

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- existence: Peano theorem
- uniqueness: assumption on y or f
- computability: assume f is computable

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If f is computable and (1) has a unique solution, then it is computable over its maximum interval of life I.

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Theorem (Buescu, Campagnolo and Graça)

Computing I (or deciding if I is bounded) is undecidable, even if f is a polynomial.

Empirical approach to complexity

Assume $y : [0, 1] \to \mathbb{R}^n$ satisfies $\forall t \in [0, 1]$:

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where $f : \mathbb{R}^n \to \mathbb{R}^n$ is Lipschitz continuous and computable.

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Useful in practice, not that much in theory.

Nonuniform complexity-theoretic approach

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But those results can be deceiving...

$$\begin{cases} y_1(0) = 1 \\ y_2(0) = 1 \\ \vdots \\ y_n(0) = 1 \end{cases} \qquad \begin{cases} y_1' = y_1 \\ y_2' = y_1 y_2 \\ \vdots \\ y_n' = y_{n-1} y_n \end{cases} \quad \rightarrow \quad y(t) = \mathcal{O}\left(e^{e^{\cdot^{\cdot^{\cdot^{e^t}}}}}\right)$$

y is PTIME over $[0,1]$

Nonuniform complexity: limitation

Example:

f PTIME analytic $\Rightarrow$ y PTIME $\Rightarrow$ $y(t) \pm 2^{-n}$ in time An^k

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$$f \text{ PTIME analytic} \;\Rightarrow\; y \text{ PTIME} \;\Rightarrow\; y(t) \pm 2^{-n} \text{ in time } An^k$$

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- “Hides” some of the complexity: A,k *could* be arbitrarily horrible depending on the dimension and f.

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But:

- “Hides” some of the complexity: A,k *could* be arbitrarily horrible depending on the dimension and f.
- Nonconstructive: might be a different algrithm for each f, or depend on uncomputable constants.

Uniform (operator) complexity approach

Assume $y : I \to \mathbb{R}^d$ satisfies $\forall t \in I$:

$$y(0) = 0, \qquad y'(t) = f(y(t)),$$

where $f : \mathbb{R}^n \to \mathbb{R}^n$ is .... Prove that $y(t) \pm 2^{-n}$ can be computed in time:

$$T(t, n, K_d, K_f)$$

where

- K_d: depends on the dimension d
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Problem: we cannot predict the behaviour of y based on f.

Parametrized complexity approach

Goal: Assume $y : I \to \mathbb{R}^d$ satisfies $\forall t \in I$:

$$y(0) = 0, \qquad\qquad y'(t) = f(y(t)),$$

where $f : \mathbb{R}^n \to \mathbb{R}^n$ is nice. Prove that $y(t) \pm 2^{-n}$ can be computed in time:

$$\text{poly}(t, n, K_d, K_f, K_y(t))$$

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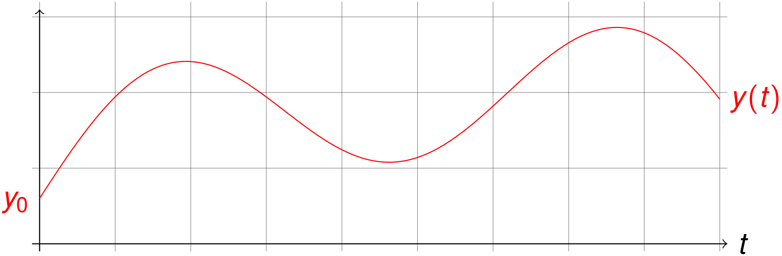
where

- K_d: depends on the dimension d
- K_f: depends on f and its representation
- K_y: is a reasonable parameter of y, ideally unknown to the algorithm (i.e. not part of the input)

Interesting parameters



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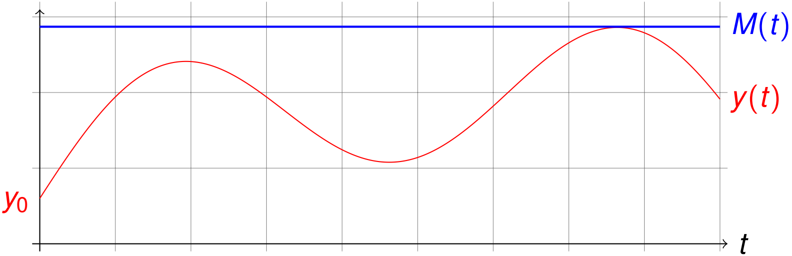

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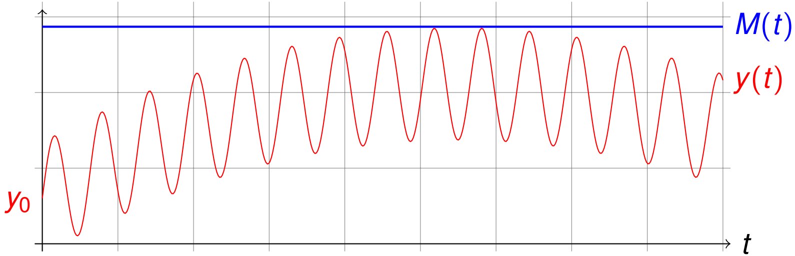

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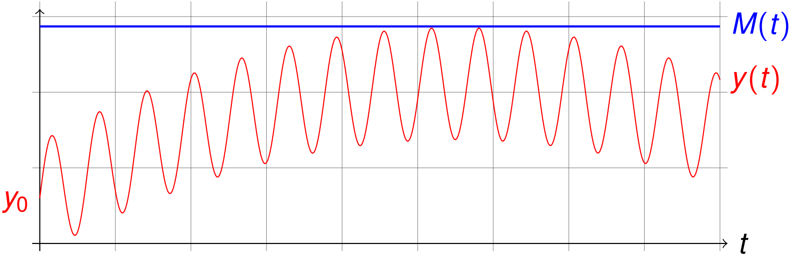

- Bounding-box: $M(t)$
- Length of the curve: $\int_0^t \|y'(u)\|\, du$

Main Result

Assume $y : I \to \mathbb{R}^d$ satisfies $\forall t \in I$:

$$y(0) = 0, \qquad y'(t) = p(y(t)),$$

where $p : \mathbb{R}^n \to \mathbb{R}^n$ is vector of multivariate polynomials.

Theorem (Our work)

Assuming $t \in I$, computing $y(t) \pm 2^{-n}$ takes time:

$$\operatorname{poly}(\deg p, \log \Sigma p, n, \ell(t_0, t))^d$$

where:

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- $\ell(t_0, t)$: "length" of y over $[t_0, t]$

$$\ell(t_0, t) = \int_0^t \max(1, \|y'(u)\|) du$$

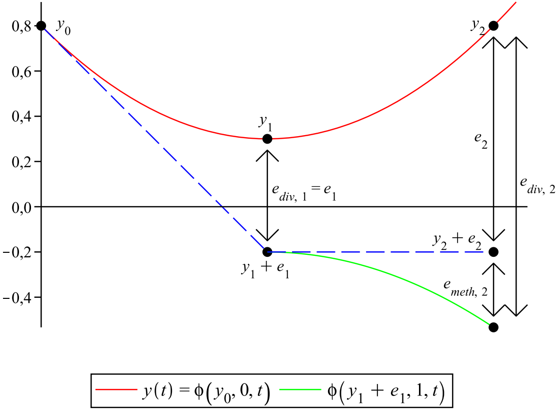
Note: the algorithm can find $\ell(0, t)$ automatically

Euler method

$$y(0) = 0 \qquad y'(t) = p(y(t)) \quad t \in I$$

Time step h, discretize compute $\tilde{y}^i \approx y(ih)$:

$$y(t+h) \approx y(t) + hy'(t) \quad \rightsquigarrow \quad \tilde{y}^{i+1} = \tilde{y}^i + hp(\tilde{y}^i)$$



Taylor method

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Lemma: $y^{(k)}(t) = P_k(y(t)) = \text{poly}(y(t))$

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$$y(t+h) \approx \sum_{j=0}^{K} \frac{h^j}{j!} y^{(j)}(t) \quad \rightsquigarrow \quad \tilde{y}^{i+1} = \sum_{j=0}^{K} \frac{h^j}{j!} P_k(\tilde{y}^i)$$

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What about h ?

- **Fixed** h: wasteful
- **Adaptive** h: choose h depending on i, p, n and $\tilde{y}^i$

Choice of the parameters

Choice of h based on an effective lower bound on radius of convergence of the Taylor series:

Lemma: If $y' = p(y)$, $\alpha = \max(1, \|y_0\|)$, $k = \deg(p)$, $M = (k-1)\Sigma p \alpha^{k-1}$ then:

$$\left\| y^{(k)}(t) - P_k(y(t)) \right\| \leqslant \frac{\alpha (Mt)^k}{1 - Mt}$$

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- local error $\approx (Mt)^k \approx 2^{-k}$: order gives the number of correct bits

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I spare you the analysis of the global error !

Conclusion

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- (Not this talk) This kind of ODE is P-complete
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Future work:

- Extend PIVP solving to more general ODEs
- Study the case when the length grows very slowly or is bounded:

$$\int_0^t \left\|y'(u)\right\| du \qquad \text{VS} \qquad \int_0^t \max(1, \left\|y'(u)\right\|) du$$

Questions ?

- Do you have any questions ?